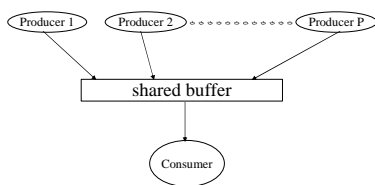


Classical IPC Problems

Problems

- producer – consumer
- readers – writers
- dining philosophers
- sleeping barber

Producer - Consumer



Producer - Consumer

- access to shared buffer through mutual exclusion
- circular buffer
- if buffer empty \rightarrow consumer waits (synchronization)

Producer – Consumer

- use counting semaphores
 - takes on ≥ 0 integers
 - used when resource capacity > 1
 - initial value = initial free resource capacity
 - P: one more unit of capacity in use
 - V: one unit of capacity freed

Producer – Consumer

- shared buffer implemented through a shared array of size N
 - array[N]
- binary semaphore:
 - mutex $\leftarrow 1$
- counting semaphores:
 - full $\leftarrow 0$: number of full buffer locations
 - empty $\leftarrow N$: number of free **buffer locations**

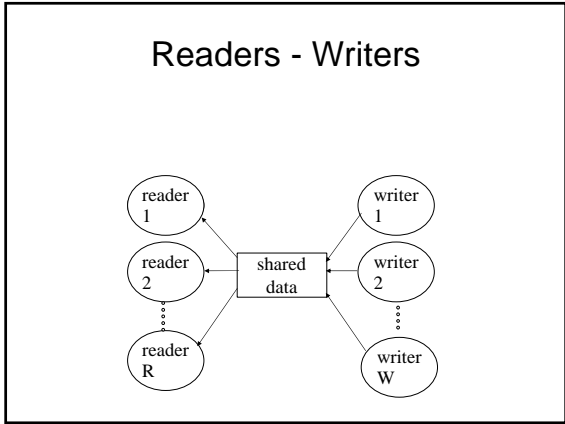
```

constant N=100;
semaphore full=0, empty=N, mutex=1;
item array[N];
int in=0, out=0;
item data;

process producer(){
while (true) {
-- produce data --
p(empty);
P(mutex);
array[in]=data;
in=(in+1)%N;
v(mutex);
v(full);
}
}

process consumer(){
while (true) {
p(full);
p(mutex);
data=array[out];
out=(out+1)%N;
v(mutex);
v(empty);
-- use data --
}
}

```



- ### Readers - Writers
- more than one reader may read shared data (no writers)
 - when a writer uses shared data, all other writers and readers must be excluded

```

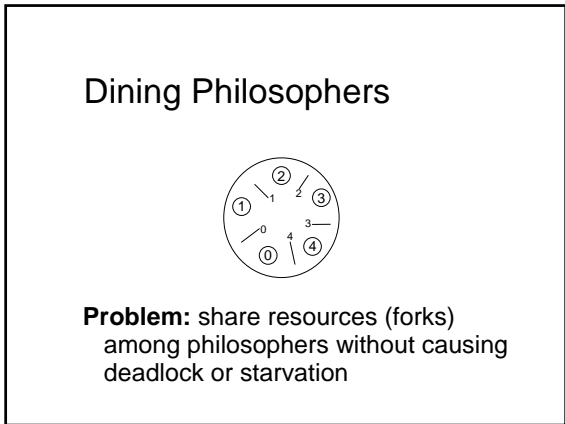
int reader=0;
semaphore read_mutex=1, data_mutex=1;

process reader() {
while (true) {
p(read_mutex);
reader=reader+1;
if (reader==1)
p(data_mutex);
v(read_mutex);
-- read data --
p(read_mutex);
reader=reader-1;
if (reader==0)
v(data_mutex);
v(read_mutex);
}
}

process writer() {
while (true) {
-- execute --
p(data_mutex);
-- write data --
v(data_mutex);
}
}

```

- ### Readers - Writers
- must find a fair solution
 - apply rules for access order:
 - if a writer is waiting for readers to be finished, do not allow any more readers
 - if a reader is waiting for a writer to finish, give reader priority



Dining Philosophers

- philosophers
 - eat pasta
 - think
- philosophers need two forks to eat

Dining Philosophers

- fact: two philosophers sitting side by side cannot eat at the same time
 - e.g. for N=5, at most 2 philosophers can eat at the same time
- solution must provide maximum amount of parallelism

Dining Philosophers

```
philosopher(i) {
  while (true) {
    think();
    take_fork(i); //left fork
    take_fork((i+4)%5); //right
    fork
    --- eat ----
    leave_fork(i);
    leave_fork ((i+4)%5);
  }
}
```

what happens if
all philosophers
take their left
forks?

Dining Philosophers

```
philosopher(i) {
  while (true) {
    think();
    take_fork(i); //left fork
    if
    (fork_free((i+4)%5)!=FALSE)
      leave_fork(i);
    else {
      take_fork((i+4)%5); //right
    }
    fork
    --- eat ----
    leave_fork(i);
    leave_fork ((i+4)%5);
  }
}
```

is it possible
that all
philosophers
starve?

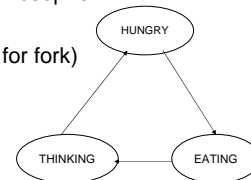
Dining Philosophers

```
philosopher(i) {
  while (true) {
    P(mutex); //binary semaphore
    think();
    take_fork(i); //left fork
    take_fork((i+4)%5); //right
    fork
    --- eat ----
    leave_fork(i);
    leave_fork ((i+4)%5);
    V(mutex);
  }
}
```

at most how
many
philosophers can
eat together?

Dining Philosophers (Correct Solution)

- state[i] : state of *i*th philosopher
 - 0 : THINKING
 - 1 : HUNGRY (wait for fork)
 - 2 : EATING



Dining Philosophers (Correct Solution)

- a philosopher can be “EATING” only if both neighbors are not “EATING”
- use a binary semaphore per philosopher
 - blocks on semaphore if a fork is not available when requested

Variables:

- N=5 philosophers
- states:
 - THINKING = 0
 - HUNGRY = 1
 - EATING = 2
- state[5]: array of size 5
- semaphores:
 - mutex ← 1
 - s[5] ← 0 array of size 5

```
process philosopher(i){
  while (true) {
    think();
    take_fork(i);
    --- eat ---
    leave_fork(i);
  }
}
```

```
leave_fork(i) {
  left=(i+1)%5;
  right=(i+4)%5;
  P(mutex);
  state[i]=THINK;
  try(left);
  try(right);
  V(mutex);
}
```

```
take_fork(i) {
  P(mutex);
  state[i]=HUNGRY;//request to eat
  try(i); //try to take forks
  V(mutex);
  P(s[i]); //blocks if can't take forks
}
```

```
try(i) {
  left=(i+1)%5;
  right=(i+4)%5;
  if ((state[i]=HUNGRY) ^
      (state[left]!=EATING) ^
      (state[right]!=EATING))
  {
    state[i]=EATING;
    v(s[i]);
  }
}
```

Sleeping Barber

- in a barber shop
 - 1 barber
 - 1 customer seat
 - N waiting seats
- barber sleeps if there are no customers
- arriving customer wakes barber up
- if barber is busy when customer arrives
 - waits if waiting seats available
 - leaves if no waiting seats available

Sleeping Barber

- 3 semaphores needed for the solution
 - customers : number of customers waiting (excluding the one in the customer seat)
 - barbers : number of available barbers (0/1 in this problem)
 - mutex : for mutual exclusion

```
constant CHAIRS=5;
int waiting=0;
semaphore customers=0,barber=0,mutex=1;
```

```
process barber() {
  while(true) {
    P(customers) ;#sleep if no customers
    P(mutex);
    waiting-- ;#remove customer
    V(barber) ;#barber ready to cut hair
    V(mutex);
    -- cut hair --
  }
}
```

```
process customer() {
  P(mutex);
  if (waiting<CHAIRS){#shop full?
    waiting++; ;#admit customer
    V(customers) ;#wake-up barber (possibly)
  }
  V(mutex);
  P(barber) ;#sleep if barber busy
  -- cut hair --
  else
    V(mutex) ;#shop is full, so leave
}
```