A WEAKLY-CONVEX FORMULATION FOR PHASELESS IMAGING

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ABSTRACT

We consider the problem of reconstructing an object given magnitudes of linear measurements. We follow the 'lifting' approach, but unlike previous work which use convex relaxations of the unit rank constraint, we use a weakly-convex matrix penalty. We derive a convergent algorithm and show that it is computationally more feasible than those obtained under convex relaxations. We demonstrate numerically that when the signal to noise ratio is high, the proposed algorithm can achieve almost error-free reconstruction with fewer measurements than when convex relaxation is employed.

Index Terms— Phaseless imaging, lifting, unit rank constraint, weakly-convex, majorization-minimization

1. INTRODUCTION

The phaseless imaging problem consists of estimating an object x given intensity measurements of the form $y_i = |\langle a_i, x \rangle|^2$, for $i = 1, 2, \ldots, m$ where a_i denotes the i^{th} row of a known imaging operator A.

A direct approach to recover x is

$$\hat{x} = \arg\min \|y - |Az|^2\|_2.$$
 (1)

Unfortunately, this formulation is non-convex and algorithms attempting to solve (1) run the risk of getting trapped in local minima [1]. An equivalent convex problem can be obtained by 'lifting' the problem [2, 3, 4]. For this, observe that $y = \text{diag}(A X A^T)$, where $X = x x^T$ and $\text{diag}(\cdot)$ extracts the diagonal of its operand. Notice that the compound operator $L(X) = \text{diag}(A X A^T)$ is linear in X. We can obtain x (upto a multiplicative factor) by estimating the unit rank positive semi-definite (psd) X via

$$\hat{X} = \arg\min_{Z \in S_n} \|y - L(Z)\|_2^2$$
 s.t.
 $\begin{cases} \operatorname{rank}(Z) = 1, \\ Z \succeq 0, \end{cases}$ (2)

where S_n denotes the set of $n \times n$ symmetric matrices. An alternative formulation is

$$\hat{X} = \arg\min_{Z \in S_n} \operatorname{rank}(Z) \quad \text{s.t.} \begin{cases} y = \operatorname{diag}(A Z A^T), \\ Z \succeq 0. \end{cases}$$
(3)

In both formulations (2), (3), the appearance of 'rank(\cdot)' makes the problem non-convex. Replacing it with a convex proxy like the nuclear norm, $\|\cdot\|_*$, one obtains a convex problem that is guaranteed to reconstruct X, under certain assumptions on A [3]. It was shown in [5] that it is also possible to drop the cost term entirely and still ensure exact recovery. Dropping the cost term leads to a feasibility problem as,

Find Z such that
$$\begin{cases} y = \operatorname{diag}(A Z A^T), \\ Z \succeq 0. \end{cases}$$
(4)

In order to obtain a simple algorithm, this problem may also be written as,

$$\min_{Z \succeq 0} \ \frac{1}{2} \| y - \operatorname{diag}(A Z A^T) \|_2^2 \tag{5}$$

For (5), the forward backward splitting algorithm [6] yields the following iterations.

1: repeat

2: $Z \leftarrow Z - \alpha L^T (L(Z) - y)$ 3: $Z \leftarrow P_+(Z)$ 4: **until** some convergence criterion is met

In the pseudocode, $P_+(\cdot)$ denotes the projection operator onto the set of positive semi definite (psd) matrices. Although this simple algorithm can correctly recover X given enough measurements, it occasionally fails to find a unit rank X, when the number of measurements are fewer. Thus, even though the unit rank constraint complicates the problem significantly by introducing non-convexity, it can be useful for eliminating incorrect estimates. In fact, the unit rank constraint allows to eliminate the psd constraint. More precisely, consider the following problem

$$\min_{Z \in S_{r}} \|y - L(Z)\|_2^2 \text{ s.t. } \operatorname{rank}(Z) = 1.$$
(6)

For this problem, we have, the following result.

Proposition 1. If \hat{X} is a local minimizer of (6), then $\hat{X} \succeq 0$.

Proof. Since X is unit rank, it can be expressed as $\hat{X} = \beta \hat{x} \hat{x}^T$, for some vector \hat{x} . Here $\beta \in \mathbb{R}$ since $\hat{X} \in S_n$. Consider the function $g(\beta) = \|y - L(\beta \hat{x} \hat{x}^*)\|_2^2$, which is a slice of the original cost, in the direction of \hat{X} . Since y > 0, the function $g(\beta)$ is differentiable. It can be shown that the derivative is strictly negative when $\beta < 0$. Thus follows the claim.

The following simple algorithm for (6) can be derived via the majorization-minimization framework [7].

2: $Z \leftarrow Z - \alpha L^T (L(Z) - y)$ \triangleright take a gradient step 3: $Z \leftarrow P_1(Z)$ \triangleright Find the closest unit rank matrix 4: **until** some convergence criterion is met

Notice that the difference between this algorithm from the former is the replacement of P_+ with P_1 , which maps Z to the unit rank matrix closest to Z. This can be achieved by computing the largest magnitude eigenvalue and the corresponding eigenvector. Thus, in general, P_1 is much easier to realize than P_+ and therefore this latter algorithm is computationally more feasible.

We have observed that the algorithm introduced above for solving (6) performs quite well but it occasionally fails to find the correct solution, even if the number of measurements is high. We think this is due to non-convex nature of the formulation (6). In order to circumvent this, we propose in this paper to relax the unit rank constraint using a penalty other than the nuclear norm and drop the psd constraint. This leads to a formulation as,

$$\min_{Z \in S_n} \frac{1}{2} \| y - L(Z) \|_2^2 + q(Z), \tag{7}$$

where a is a penalty function on Z, derived from our previous work [8]. The parameters of q can be chosen such that q(Z) is arbitrarily small for a unit rank Z and increases rapidly if more than one singular value of Z is non-zero. This property allows to significantly reduce a possible bias introduced by the addition of $q(\cdot)$, provided that a unit rank solution is found for the problem. The proposed qis weakly-convex, that is, q becomes convex after adding a quadratic function [9]. Thus, in contrast to earlier work, the non-convexity introducing rank constraint is not discarded [5] or relaxed to a nuclear norm [2, 3], but relaxed using a function that provides a better approximation of the constraint. The algorithm for solving (7) can be obtained by replacing P_1 with the proximity operator of q [6] and will be derived in the sequel. We will show that the proximity operator of q can be realized almost as easily as P_1 . Therefore, the computational price paid for the proposed relaxation is lower, compared to the convex relaxations mentioned above.

Other Related Work : In optical imaging, the phaseless imaging problem described above is also known as the phase retrieval problem [1, 4]. Simple alternating projection algorithms, without convex relaxation have been described in [1]. However, these algorithms do not come with theoretical convergence guarantees and in practice it has been observed that they can get stuck at non-optimal solutions. For the phase retrieval problem, the authors of [10] proposes to reconstruct the phases of the measurements and after a lifting step followed by discarding of a unit rank constraint, they obtain a convex semi-definite programming problem. More recently, [11] proposes to iteratively reconstruct the support of x, without lifting the problem to a higher dimension. At each iteration, the support is updated based on the entries of the gradient of the cost function in (1). Another algorithm that does not rely on lifting is presented in [12], based on generalized approximate message passing. These latter approaches are attractive in that they do not need to introduce new variables. However, they cannot offer global convergence guarantees due to the inherent non-convexity of the problem. An alternative approach, which also does not rely on lifting, is proposed in [13]. Starting with a good initialization, the algorithm applies a form of gradient descent in order to minimize a quadratic loss function. Assuming that the measurement matrix is obtained by sampling a certain probability distribution, the authors show that, with high probability, the proposed algorithm converges to a minimizer of the cost function.

As briefly noted above, our idea is to improve the lifting scheme by relaxing the non-convex problem to a weakly-convex problem, instead of a convex problem as is done in previous work. While weakly-convex problems are also non-convex, non-convexity is reduced, in some sense. Our numerical experiments suggest that, besides reduced computational load, weakly-convex relaxation, in comparison to convex relation, improves performance and yields a sharper transition in the phase diagram between success and failure of reconstruction.

Notation : Throughout the text, we work with real variables for simplicity. The subscript T denotes transpose. Most results carry over to the complex case straightforwardly, simply by replacing T with the Hermitian conjugate * , with the exception of Lemma 1, which requires a more careful modification.

Outline : We introduce the penalty and discuss its properties in Section 2. We derive an algorithm for (7) and show that it converges

in Section 3. We provide numerical experiments demonstrating the performance of the proposed algorithm as well those mentioned in the introduction in Section 4, followed by concluding remarks in Section 5.

2. THE PENALTY FUNCTION AND THE ALGORITHM

We will derive the penalty function $q : S_n \to \mathbb{R}_+$ by adapting a penalty proposed earlier for \mathbb{R}^n [8].

2.1. A Penalty on \mathbb{R}^n

Consider the following function defined on \mathbb{R}^n .

$$p_{\lambda,\gamma}(x) = \lambda \bigg[\|x\|_1 + \gamma \bigg(\sum_{i=1}^{n-1} \sum_{j=i+1}^n |x_i \, x_j| \bigg) \bigg].$$
(8)

Observe that if $||x||_0 \leq 1$, then the term enclosed in parentheses vanishes. The proximity operator [6] for this function is defined as,

$$J_{\alpha p}(x) = \arg \min_{z} \frac{1}{2\alpha} ||x - z||_{2}^{2} + p_{\lambda,\gamma}(z).$$
(9)

If $\alpha \lambda \gamma < 1$, then the cost function in (9) is strictly convex so that $J_{\alpha p}$ is well-defined. This follows from the observation that the term in (8) enclosed in parentheses can be written as $(||x||_1^2 - ||x||_2^2)/2$. To describe $J_{\alpha p}$, suppose $z = J_{\alpha p}(x)$, and assume (for notational convenience) that $|x_i|$'s are ordered (i.e., $|x_i| \ge |x_{i+1}|$). If k components of z are non-zero, then z and x are related as,

$$z_i = (1 - \alpha \lambda \gamma)^{-1} \operatorname{soft}(x_i, \tau_k), \qquad (10)$$

where the soft threshold operator is defined as $soft(a, \tau) = sign(a) \max(|a|-\tau, 0)$ and τ_k is a signal dependent threshold given as,

$$\tau_k = \frac{\alpha \lambda \left(1 - \alpha \lambda \gamma\right) + \alpha \lambda \gamma \sum_{i=1}^k |x_i|}{1 + (k-1) \alpha \lambda \gamma}.$$
 (11)

We remark that the description above is implicit since the number of non-zeros of z (i.e., the integer k) is not known in advance. However, a candidate \tilde{k} can be validated by checking whether the threshold $\tau_{\tilde{k}}$ produces a z with \tilde{k} non-zeros [8]. Therefore, a linear search can be performed starting from $\tilde{k} = 1$ to find the correct value of k. Here we note a special case of interest, which can also be found in [8].

Proposition 2. For $x, z \in \mathbb{R}^n$, suppose $z = J_{\alpha p}(x)$. Suppose also that x_l denotes the entry of x with the largest magnitude. If

$$\alpha \,\lambda \,\gamma > \max_{i \neq l} \frac{|x_i| - \alpha \,\lambda}{|x_l| - \alpha \,\lambda} \tag{12}$$

then $z_l = \operatorname{soft}(x_l, \alpha \lambda)$ and $z_i = 0$, for $i \neq l$.

In words, all but the largest (in magnitude) component of x is kept if the ratio of the largest component to the rest is greater than some threshold. This property is the key to imposing the unit rank constraint, as will be clarified below.

2.2. The Induced Penalty on S_n

We extend the domain of the penalty p from \mathbb{R}^n to S_n as follows. Suppose the eigendecomposition of $X \in S_n$ is given as $X = V \Lambda V^T$. We define

$$q_{\lambda,\gamma}(X) = p_{\lambda,\gamma}(\operatorname{diag}(\Lambda)). \tag{13}$$

The proximity operator for q is defined as,

$$J_{\alpha q}(X) = \arg \min_{Z \in S_n} \frac{1}{2\alpha} \|X - Z\|_F^2 + q_{\lambda,\gamma}(Z).$$
(14)

It follows by the discussion in [14] (see specifically [15]) along with the properties of p that $J_{\alpha q}$ is well-defined when $\alpha \lambda \gamma < 1$. This proximity operator is related to $J_{\alpha p}$ as follows [15].

Proposition 3. Suppose the eigendecomposition of $X \in S_n$ is $X = V \Lambda V^T$. Also, let Λ' denote the diagonal matrix that satisfies diag $(\Lambda') = J_{\alpha p}(\operatorname{diag}(\Lambda))$. Then, $J_{\alpha q}(X) = V \Lambda' V^T$.

Suppose $Z = J_{\alpha q}(X)$. If Z has k non-zero eigenvalues, then it follows by the discussion preceding Prop. 2 that we need to compute only the largest magnitude k + 1 eigenvalue/eigenvector pairs of X. For k = 1 (the unit rank Z case), we have the following result, which is a corollary of Prop. 2 and Prop. 3.

Corollary 1. Suppose the eigendecomposition of $X \in S_n$ is X = $V \Lambda V^T$ and Λ_l is the diagonal entry of Λ with the largest magnitude. If

$$\alpha \lambda \gamma > \max_{i \neq l} \frac{\left(|\Lambda_i| - \alpha \lambda\right)_+}{\left(|\Lambda_l| - \alpha \lambda\right)_+} \tag{15}$$

then, $J_{\alpha q}(X) = \operatorname{soft}(\Lambda_l, \alpha \lambda) v_l v_l^T$, where v_l is the l^{th} column of

This corollary suggests that if the gap between the largest eigenvalue and the rest of the eigenvalues of X, is large enough, $J_{\alpha q}(\cdot)$ reduces X to a unit rank matrix. We next derive and study the convergence of a monotone descent algorithm for solving (7).

3. A MINIMIZATION ALGORITHM

Recall the formulation in (7).

$$\min_{Z \in S_n} \left\{ C(Z) = \frac{1}{2} \left\| y - L(Z) \right\|_2^2 + q_{\lambda,\gamma}(Z) \right\},$$
(16)

It can be shown that provided $L^T L \succ (\lambda \gamma)I$, this problem is convex. However, even if this is not satisfied, it is possible to derive a monotone descent algorithm by employing majorizationminimization techniques [16, 7].

3.1. Derivation of the Algorithm

Suppose Z^k denotes our k^{th} iterate. Consider the following update.

$$Z^{k+1} = \arg \min_{Z \in S_n} \left\{ C_k(Z) = C(Z) + \frac{1}{2\alpha} \|Z - Z^k\|_F^2 - \frac{1}{2} \|L(Z - Z^k)\|_2^2 \right\}.$$
 (17)

Observe that $C_k(Z^k) = C(Z^k)$. Provided $\alpha \sigma(L^T L) < 1$, it follows that C_k upper bounds C, that is, $C_k(Z) \ge C(Z)$ for all $Z \in S_n$. In this case, we will have $C(Z^{k+1}) \le C(Z^k)$, where the inequality is strict if $Z^k \neq Z^{k+1}$. To find an expression for Z^{k+1} , we expand C_k and find that

$$Z^{k+1} = \arg\min_{Z \in S_n} \frac{1}{2\alpha} \| Z - \bar{Z}^k \|_F^2 + q_{\lambda,\gamma}(Z),$$
(18)

where $\bar{Z}^k = Z^k - \alpha L^T (L(Z^k) - y)$. Provided that $\alpha < 1/(\lambda \gamma)$, we have

$$Z^{k+1} = J_{\alpha q} \left(\bar{Z}^k \right). \tag{19}$$

Algorithm 1 A Descent Algorithm for (7) 1: $k \leftarrow 1$, initialize Z^1 , set α (see Prop. 4) 2: repeat $Z \leftarrow Z^k - \alpha \, L^T \left(L(Z^k) - y \right)$ ⊳ take a gradient step 3: $Z^{k+1} \leftarrow J_{\alpha q}(Z)$ 4: \triangleright Apply the prox. operator for q $k \leftarrow k+1$ 5: 6: until some convergence criterion is met

A summary is provided in Algorithm 1.

The foregoing discussion, along with Thm. 3.1 of [17], leads to the following proposition.

Proposition 4. Suppose $\alpha < \min(1/(\lambda, \gamma), 1/\sigma(L^T L))$. Then, Z^{k} 's in Algorithm 1 monotonically decrease the cost, i.e., $C(Z^{k+1}) \leq C(Z^{k})$ and the sequence of cost values $c_k = C(Z^{k})$ converge to a value $C(Z^*)$, where Z^* is a stationary point of $C(\cdot)$.

3.2. Convergence of the Iterates

In practice, we are not only interested in the convergence of the sequence of cost values. It is also desired that the sequence of iterates themselves converge. This is addressed in the following.

Proposition 5. Suppose $\alpha < \min(1/(\lambda \gamma), 1/\sigma(L^T L))$. Then, Z^k 's in Algorithm 1 converge to a local minimizer of (7).

Prop. 5 is a consequence of Thm. 5.1 of [18] which implies that the iterates of monotone descent algorithms converge to a local minimum when the cost function is semi-algebraic¹. Sums of semialgebraic functions are semi-algebraic [18] and a quadratic function (like the first term of $C(\cdot)$) is semi-algebraic since it is a polynomial of its arguments. Therefore, for Prop. 5, it is sufficient to show that q is semi-algebraic, which is addressed in the following lemma.

Lemma 1. Viewing an $n \times n$ matrix X as an element of \mathbb{R}^{n^2} , the nuclear norm $||X||_*$ and the regularizer q introduced in (13) are semialgebraic functions defined on \mathbb{R}^{n^2} .

Proof. We start by showing that the nuclear norm is semi-algebraic. Let Q denote the subset of \mathbb{R}^{n^2} associated with orthogonal matrices (matrices are formed by placing the components in a specific order). Notice that \mathcal{Q} (or $\mathcal{Q} \times \mathcal{Q}$) is a compact semi-algebraic set, since it is the solution set of a finite number of polynomial equations. Observe also that for $Q_i \in Q$, we have $\operatorname{tr}(Q_1^T X Q_2) \leq ||X||_*$, where equality is satisfied if the SVD of X is $X = Q_1 \Sigma Q_2^T$. Therefore,

$$||X||_* = \max_{Q_1, Q_2 \in \mathcal{Q} \times \mathcal{Q}} \{h(X, Q_1, Q_2) = \operatorname{tr}(Q_1 X Q_2)\}.$$
 (20)

But in this equation $h(X, Q_1, Q_2)$ is a polynomial of its arguments. Thus the function $||X||_*$ is semi-algebraic (see e.g. the discussion following Thm 2.2 in [18]). Notice now that q(X) = $||X||_*^2 - ||X||_2^2$. Since the sum and products of semi-algebraic functions are semi-algebraic, it follows that q is semi-algebraic.

¹In a nutshell, a function is semi-algebraic if its graph can be expressed as a finite intersection of polynomial inequalities. Due to lack of space, we refer to [18] and the references therein for the precise definition and discussion of semi-algebraic sets and functions.



Fig. 1. The average recovery error for different signal lengths and number of measurements for (a) optimizationless reconstruction [5], (b) unit rank algorithm, (c) the proposed reconstruction algorithm. (d) shows slices of E_x for all three algorithms for n = 50.

4. NUMERICAL EXPERIMENTS

In this section, we compare the proposed algorithm with the two algorithms provided in the introduction that solve (5) and (6), which we will refer to as 'optimizationless'² and 'unit rank' in what follows. For the purpose of comparison, we performed experiments similar to those of [5]. We remark that the optimizationless algorithm has been used earlier in [10] as a benchmark algorithm (see, in particular, comparisons with PhaseLift [3]). We also note that, at least in the high SNR case, the trace term in the PhaseLift formulation introduces a certain bias and the reconstruction is different than those obtained by the algorithms considered in the experiments below, which are almost error-free, provided that the number of observations are sufficient (see Fig. 3 in [5]).

Experiment 1. We randomly produce the object x by uniformly sampling the *n*-dimensional sphere, and the $m \times n$ observation matrix, A, by sampling a Gaussian distribution. We produce the observations as,

$$y = |Ax|^2 + u,$$
 (21)

where u denotes a Gaussian noise term. We can define two different recovery errors, related to x and $X = x x^T$. If \hat{X} denotes the reconstruction obtained by an algorithm, the normalized error for recovering X can be taken as $E_X = \|\hat{X} - X\|_F / \|X\|_F$. Given the matrix estimate X, we take $\hat{x} = \sqrt{\lambda} e$, where λ is the greatest eigenvalue of \hat{X} and e is the corresponding unit norm eigenvector. The recovery error for x (recall that x is unit norm) is then defined as $E_x = \min_{k=0,1} \|(-1)^k \hat{x} - x\|_2$. We remark that [5] uses E_X but we think that E_x is of interest since x is the object we aim to reconstruct. We have not seen a marked difference between the behaviors of the two distinct error terms. In the setup described above, for each

Recovery Error wrt. Number of Measurements



Fig. 2. Reconstruction error for varying number of measurements, for a signal of length 100. (a) SNR = 50 dB, (b) SNR = 5 dB.

m and n value, we average the recovery errors over twenty trials in order to obtain an average figure. We set the constants used in the algorithms as $\alpha = 10^{-4}$, $\lambda = 10^{-5}$, $\gamma = 0.05/\lambda$. Notice that, with these settings, we have $\alpha \lambda \gamma < 1$, as required in the proposed algorithm.

Figure 1 shows the phase transition diagrams as n and m vary, where the observation SNR is 50 dB. To produce these figures, we ran the proposed algorithms for 1000 iterations. We observe that the proposed algorithm and the unit rank algorithm both have sharper transitions between success and failure, compared to optimization-less reconstruction. Compared to unit rank, the proposed algorithm performs more stably and requires fewer measurements for successful recovery. In particular, we see in Fig.1d that unit rank has not improved much when the number of measurements, m, is in the range [100, 150] and has an unexpected local maximum at m = 175, whereas the proposed algorithm steadily improves as m increases.

Experiment 2. In order to demonstrate the sharp transition noted above and test robustness against noise, we performed similar experiments for n = 100 with different SNR values. In this case, we set $\alpha = 10^{-4}/5$ (to ensure convergence) where the other parameters are kept as in the previous experiment. The reconstruction errors E_x are shown in Fig. 2 for SNR values 50 dB and 5 dB. In this case, we ran the algorithms for 5000 iterations, to avoid premature stopping.

In both cases, the average error for optimizationless recovery appears to be smooth with respect to m and especially towards the right end of the graphs, it exceeds the error obtained by the proposed algorithm. As for the unit rank algorithm and the proposed algorithm, we see a similar trend as in the previous experiment. The proposed algorithm achieves a lower error than unit rank for all m. We also observe that the proposed algorithm monotonically decreases the average error as m increases (except when the number of measurements is too few), whereas the unit rank algorithm can produce incorrect estimates even when $m \geq 300$ – e.g., when m = 400 in Fig. 2b.

5. CONCLUSION

The proposed algorithm falls somewhere in between algorithms obtained via convex relaxation and an algorithm that imposes the nonconvex unit rank constraint at each iteration, both in terms of performance as well as computational load. Therefore, *weak*-convex relaxation allows to achieve a 'tunable' performance. Currently, due to lifting, the number of variables are significantly increased and hence memory requirements are high for real-life problems. In order to make the approach feasible for larger problems, we plan to complement the algorithm with a smart initialization strategy to achieve a reasonable approximation with a fewer number of iterations.

²For 'optimizationless' reconstruction, we use the accelerated version of the algorithm as described in [5].

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