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Buffer Management

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A journey of a byte

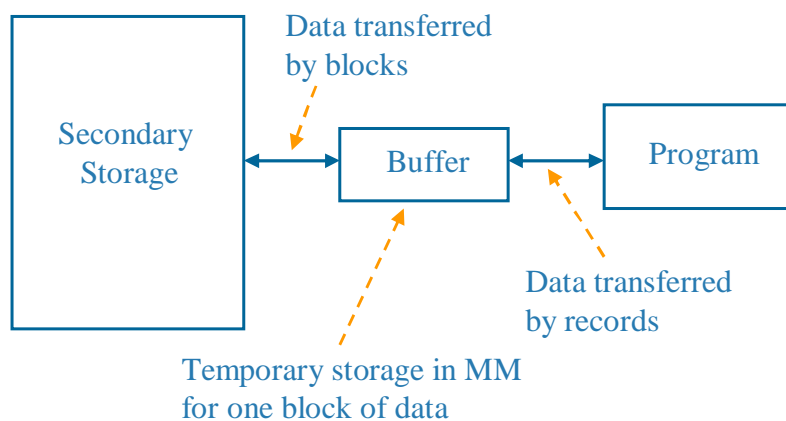
- ▶ Suppose in our program we wrote:
`outfile << c;`
- ▶ This causes a call to the **file manager** (a part of O.S. responsible for I/O operations)
- ▶ The O/S (File manager) makes sure that the byte is written to the disk.
- ▶ Pieces of software/hardware involved in I/O:
 - Application Program
 - Operating System/ file manager
 - I/O Processor
 - Disk Controller

- ▶ **Application program**
 - Requests the I/O operation
- ▶ **Operating system / file manager**
 - Keeps tables for all opened files
 - Brings appropriate sector to buffer.
 - Writes byte to buffer
 - Gives instruction to I/O processor to write data from this buffer into correct place in disk.
 - Note: the buffer is an exact image of a cluster in disk.
- ▶ **I/O Processor**
 - a separate chip; runs independently of CPU
 - Find a time when drive is available to receive data and put data in proper format for the disk
 - Sends data to disk controller
- ▶ **Disk controller**
 - A separate chip; instructs the drive to move R/W head
 - Sends the byte to the surface when the proper sector comes under R/W head.

Buffer Management

- ▶ Buffering means working with large chunks of data in main memory so the number of accesses to secondary storage is reduced.
- ▶ Today, we'll discuss the System I/O buffers. These are beyond the control of application programs and are manipulated by the O.S.
- ▶ Note that the application program may implement its own "buffer" – i.e. a place in memory (variable, object) that accumulates large chunks of data to be later written to disk as a chunk.

System I/O Buffer



Buffer Bottlenecks

- ▶ Consider the following program segment:

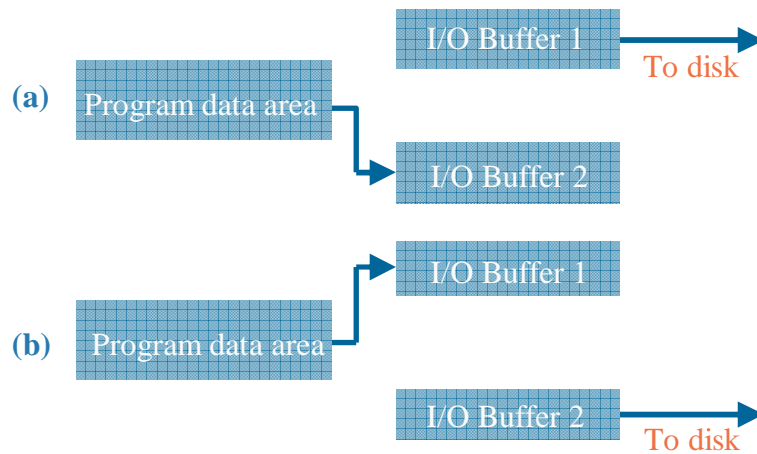
```
while (1) {
  infile >> ch;
  if (infile.fail()) break;
  outfile << ch;
}
```

- ▶ What happens if the O.S. used only one I/O buffer?
⇒ Buffer bottleneck
- ▶ Most O.S. have an input buffer and an output buffer.

Buffering Strategies

- ▶ **Double Buffering:** Two buffers can be used to allow processing and I/O to overlap.
 - Suppose that a program is only writing to a disk.
 - CPU wants to fill a buffer at the same time that I/O is being performed.
 - If two buffers are used and I/O-CPU overlapping is permitted, CPU can be filling one buffer while the other buffer is being transmitted to disk.
 - When both tasks are finished, the roles of the buffers can be exchanged.
- ▶ The actual management is done by the O.S.

Double Buffering



Other Buffering Strategies

- ▶ **Multiple Buffering:** instead of two buffers any number of buffers can be used to allow processing and I/O to overlap.
- ▶ **Buffer pooling:**
 - There is a pool of buffers.
 - When a request for a sector is received, O.S. first looks to see that sector is in some buffer.
 - If not there, it brings the sector to some free buffer. If no free buffer exists, it must choose an occupied buffer. (usually LRU strategy is used)

Buffering Strategies: Move & Locate mode

- ▶ Move mode (using both system buffer & program buffer)
 - moving data from one place in RAM to another before they can be accessed
 - sometimes, unnecessary data moves
- ▶ Locate mode (using system buffer only or program buffer only)
 - perform I/O directly between secondary storage and program buffer (program's data area)
 - system buffers handle all I/Os, but program uses locations through pointer variable

Move Mode and Location Mode

