

## Nature-Inspired Computing

### Grammatical Evolution

Dr. Şima Uyar  
September 2007

## Overview

- Michael O'Neill and Conor Ryan
- 1998 – 2001
- binary chromosomes
- grammar defined through BNF
- genotype to phenotype mapping
- regular genetic algorithm operators
- produces syntactically correct programs in any language (originally developed for Lisp)

## Backus Naur Form (BNF)

- notation for expressing the grammar of a language in the form of production rules
- BNF grammars consist of:
  - terminals (items which appear in language, e.g. + or -)
  - non-terminals (can be expanded into one or more terminals or non-terminals)

## Backus Naur Form (BNF)

- a grammar represented by tuple  $\{N, T, P, S\}$ 
  - N: set of non-terminals
  - T: set of terminals
  - P: set of production rules mapping elements of N to T
  - S: a start symbol which is a member of N

## Backus Naur Form (BNF)

- when more than one production can be applied to one element of N, they are separated by the "|" symbol

## Backus Naur Form (BNF)

### example:

```
N={expr,op, pre-op,var}
T={sin,cos,+,-,/,*,X,1.0,(,)}
S=expr
```

## Backus Naur Form (BNF)

### example (contd.):

P can be represented as:

(A) <expr> ::= <expr> <op> <expr>	(a)
(<expr> <op> <expr>)	(b)
<pre-op> (<expr>)	(c)
<var>	(d)
(B) <op> ::= +	
(e)	
-	(f)
/	(g)
*	(h)
(C) <pre-op> ::= sin	(i)
cos	(j)
(D) <var> ::= X	(k)
1.0	(l)

## Grammatical Evolution

- 8 bits = codons
- integer values of codons used to select production rule
- wrapping of chromosome is used

## Grammatical Evolution

- beginning from left hand side of genome, codon integer values are generated and used to select rules from BNF grammar until one of these arise:
  - a complete program is generated (all non-terminals in expression is transformed into elements from the terminal set)
  - end of genome reached; wrapping operator is invoked; continue from left; repeat until max no of wrapping reached
    - if max no of wrappings reached and individual is not complete, it is assigned the lowest possible fitness

## Example Individual

220 40 16 203 101 53 202 203 102 55 220  
202 19 130 37 202 203 32 39 202 203 102

Step 1: S=expr ⇒ 4 to choose from in rule A  
220 mod 4 = 0, so first production is chosen  
<expr> replaced with <expr><op><expr>

## Example Individual

220 40 16 203 101 53 202 203 102 55 220  
202 19 130 37 202 203 32 39 202 203 102

Step 2: continue with first <expr>: 40 mod 4=0, so  
<expr><op><expr><op><expr>

Step 3: 16 mod 4 = 0 so  
<expr><op><expr><op><expr><op><expr>  
Step 4: 203 mod 4=3 so  
<var><op><expr><op><expr><op><expr>  
Step 5: 101 mod 2 =1  
<1.0><op><expr><op><expr><op><expr>  
Step 6: 53 mod 4=1  
<1.0><-><expr><op><expr><op><expr>  
Step 6: 202 mod 4=2  
<1.0><-><pre-op>(<expr>)<op><expr><op><expr>  
Step 7: only one choice for pre-op, no codon used  
<1.0><-><sin>(<expr>)<op><expr><op><expr>  
Step 8: 203 mod 4=3  
<1.0><-><sin>(<var>)<op><expr><op><expr>  
Step 9: 102 mod 2=0  
1.0-sin(x)<op><expr><op><expr>  
...  
...  
Outcome: 1.0-sin(x)\*sin(x)-sin(x)\*sin(x)

## Grammatical Evolution

- issues:
  - degeneracy (due to modulo operator)
  - neutral mutations
  - disruptive effect of one point crossover
    - work on special crossovers
- many applications
- work still continues